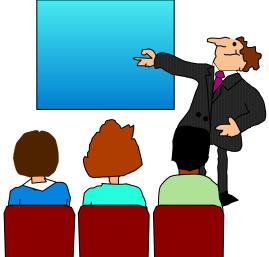
**Assembler Language "Boot Camp" Part 5 - Decimal and Logical Instructions** SHARE in San Francisco August 18 - 23, 2002 Session 8185



- Who are we?
  - John Dravnieks, IBM Australia
  - John Ehrman, IBM Silicon Valley Lab
  - Michael Stack, Department of Computer Science, Northern Illinois University

#### Who are you?

- An applications programmer who needs to write something in S/390 assembler?
- An applications programmer who wants to understand S/390 architecture so as to better understand how HLL programs work?
- A manager who needs to have a general understanding of assembler?

Our goal is to provide for professionals an introduction to the S/390 assembly language

These sessions are based on notes from a course in assembler language at Northern Illinois University

The notes are in turn based on the textbook, <u>Assembler Language with ASSIST and</u> <u>ASSIST/I</u> by Ross A Overbeek and W E Singletary, Fourth Edition, published by Macmillan

- The original ASSIST (<u>Assembler System for</u> <u>Student Instruction and Systems Teaching</u>) was written by John Mashey at Penn State University
- ASSIST/I, the PC version of ASSIST, was written by Bob Baker, Terry Disz and John McCharen at Northern Illinois University

- Both ASSIST and ASSIST/I are in the public domain, and are compatible with the System/370 architecture of about 1975 (fine for beginners)
- Both ASSIST and ASSIST/I are available at http://www.cs.niu.edu/~mstack/assist

Other materials described in these sessions can be found at the same site, at http://www.cs.niu.edu/~mstack/share

Please keep in mind that ASSIST and ASSIST/I are not supported by Penn State, NIU, or any of us

Other references used in the course at NIU:

- Principles of Operation
- System/370 Reference Summary
- High Level Assembler Language Reference

Access to PoO and HLASM Ref is normally online at the IBM publications web site

Students use the S/370 "green card" booklet all the time, including during examinations (SA22-7209)

### **Our Agenda for the Week**

Session 8181: Numbers and Basic Arithmetic

Session 8182: Instructions and Addressing

Session 8183: Assembly and Execution; Branching

### **Our Agenda for the Week**

Session 8184: Arithmetic; Program Structures

Session 8185: Decimal and Logical Instructions

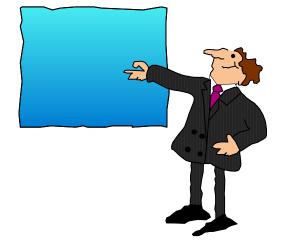
Session 8186: Assembler Lab Using ASSIST/I



- The SI and SS Instruction Formats
- Decimal Arithmetic
- Instructions for Logical Operations



## The SI and SS Instruction Formats





This format encodes the second operand as an "immediate" data byte within the instruction

The symbolic instruction format is

label mnemonic address, byte

The encoded form of an SI instruction is

 $\mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{I2}} \mathbf{h}_{\mathrm{I2}} \mathbf{h}_{\mathrm{B1}} \mathbf{h}_{\mathrm{D1}} \mathbf{$ 

MOVE IMMEDIATE is our first SI instruction

- **label** MVI  $D_1(B_1), I_2$
- Stores a copy of the immediate byte, I<sub>2</sub>, at the memory location given by D<sub>1</sub>(B<sub>1</sub>)

- The second operand can be specified as a decimal number or as any one-byte value valid in DC
  - Decimal: 91
  - Hexadecimal: X ' 5B '
  - Binary: **B'01011011'**
  - Character: C'\$'

# For example, to place a single blank at PLINE

MVI PLINE,C''

- The COMPARE LOGICAL IMMEDIATE instruction compares the byte in memory to the immediate data byte as unsigned binary numbers
  - **label** CLI  $D_1(B_1), I_2$

CLI sets the condition code in the same way as other compare instructions

The following code sample scans an 80-byte data area labelled CARD and replaces blanks with zeros

	•••		
	LA	4,CARD	Start scan here
	LA	3,80	and scan 80 bytes
SCAN	CLI	0(4),C''	Look for blank
	BNE	BUMP	Branch if not blank
	MVI	0(4),C'0'	Else change to 0
BUMP	LA	4,1(,4)	Move to next byte
	BCT	3, SCAN	Continue for 80

- In this format, which occupies 3 halfwords, both operands reference memory locations, and there is either one 256-byte-max length or two 16-byte-max lengths
- The symbolic instruction format is either
  - label mnemonic addr1(len),addr2 or
  - label mnemonic addr1(len1),addr2(len2)

- Each SS instruction is defined to have one of the length formats; we will see only the first for now
- The encoded form of an SS instruction is
  - $\mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{L}} \mathbf{h}_{\mathrm{L}} \quad \mathbf{h}_{\mathrm{B1}} \mathbf{h}_{\mathrm{D1}} \mathbf{h}_{\mathrm{D1}} \mathbf{h}_{\mathrm{D1}} \quad \mathbf{h}_{\mathrm{B2}} \mathbf{h}_{\mathrm{D2}} \mathbf{h}_{\mathrm{D2}} \mathbf{h}_{\mathrm{D2}} \mathbf{h}_{\mathrm{D2}} \quad \mathrm{Or}$
  - $\mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{OP}} \mathbf{h}_{\mathrm{L1}} \mathbf{h}_{\mathrm{L2}} \mathbf{h}_{\mathrm{B1}} \mathbf{h}_{\mathrm{D1}} \mathbf{h}_{\mathrm{D1}} \mathbf{h}_{\mathrm{D1}} \mathbf{h}_{\mathrm{B2}} \mathbf{h}_{\mathrm{D2}} \mathbf{$

- Very Important: the <u>encoded length</u> is one less than the <u>symbolic length</u> (as well as the <u>effective length</u>) and is referred to as the "<u>length code</u>"
- Thus, in the first format, 1 to 256 bytes may be specified where 0 to 255 is encoded
- An explicit length of 0 results in an encoded length of 0, so the effective length is 1

# MOVE CHARACTERS is our first SS instruction

- **label** MVC  $D_1(L,B_1), D_2(B_2)$
- Copies from 1 to 256 bytes from the second operand location to the first

- For example, to copy 8 bytes from the location addressed by register 1 to 14 bytes beyond the location addressed by register 12
  - Symbolic: MVC 14(8,12),0(1)
  - Encoded: D207 C00E 1000
  - Note the encoded length!

- Implicit addresses may be used, of course, and with or without an explicit length
  - MVC FIELD1(15),FIELD2
  - MVC FIELD1,FIELD2

Both generate the same object code if FIELD1 (the first operand) has a "length attribute" of 15, as in

FIELD1 DS CL15

- Any explicit length will take precedence over the implicit length derived from the length attribute
- So, in the previous example the following instruction will move only 8 bytes, even though **FIELD1** has a length of 15
  - MVC FIELD1(8),FIELD2
- Implicit lengths change automatically when data lengths changes

The effect of MVC is to replace L bytes beginning at the first operand location with a copy of the L bytes beginning at the second operand location

The target is altered, one byte at a time, starting on the "left" (the beginning, or low, address)

- This means that the fields can overlap with predictable results, and here is an historically important example
- There is often a "print buffer" in which output lines are constructed, and after printing a line, the buffer should be cleared to blanks (this assumes the PLINE has a length attribute of 133, as it would if PLINE DS CL133 is used)

So, we would normally clear the buffer by copying a string of blanks to it
 MVC PLINE,=CL133' '

But by using the overlap, we can "save" 129 bytes

- MVI PLINE,C''
- MVC PLINE+1(132),PLINE

#### Suppose FIELD DC C'123456' What is FIELD after MVC FIELD+2(4),FIELD ?

- Another SS instruction which uses the first length format is COMPARE LOGICAL CHARACTERS
  - **label** CLC  $D_1(L,B_1), D_2(B_2)$
- As with all compares, this just sets the condition code
- The operation stops when the first unequal bytes are compared

## **Decimal Arithmetic**

In Which We Switch to Counting on Our Fingers or Toes Instead of Our Hands

2

Thus far, the computations we've done have been with binary data

This is not always satisfactory, especially when financial calculations are required

For example, decimal percentages are inaccurate in binary (try long division on  $1/10_{10} = 1/1010_2 = .000110011...$ )

This (infinite repetition) annoys auditors

The solution is to use computers with decimal data types and instructions

There are two decimal data formats

- Zoned Decimal associated with I/O operations
- Packed Decimal used for decimal arithmetic

- A <u>zoned decimal</u> number is a sequence of bytes in which each byte has
  - a decimal digit 0-9 in the right digit and
  - a zone digit (hex F) in the left digit, except that the rightmost zone is the sign

- That is, a zoned decimal number has the format
  - ZdZdZd...sd where
    - Z is the hex digit F
    - •d is a decimal digit 0-9
    - **s** is the sign
      - C, A, F, or E means + (C is preferred)
      - D or B means (D is preferred)

An example is F1F2C3, for +123

- A zoned number is very close to the EBCDIC representation of its value, except that the rightmost byte has a sign, so won't print as a number
  - So our zoned +123 prints as **12C**

- A packed decimal number has the zones removed and the sign switched with its digit; that is,
  - dddddd...ds
- Note that there are always an odd number of digit positions in a packed decimal number

The assembler can generate data of types Z (zoned) and P (packed)

#### **Decimal Data**

- 📃 label DC mZLn'z'
  - DC Z'+123' = F1F2C3
  - **DC** ZL3'-1.2' = F0F1D2
- label DC mPLn'p'
  - DC P'+123' = 123C
  - DC P'-1.2' = 012D
  - DC PL2'1234' = 234C

Both of these are SS instructions of the second type - that is, <u>each operand has a</u> <u>length field</u> which will accommodate a length code of 0-15 (so the effective length is 1-16 bytes)

- Use the PACK instruction to convert a number from zoned decimal to packed decimal
- Use the UNPK instruction to convert a number from packed decimal to zoned decimal

#### 

- The rightmost byte of the second operand is placed in the rightmost byte of the first operand, with zone (sign) and numeric digits reversed
- The remaining numeric digits from operand 2 are moved to operand 1, right to left, filling with zeros or ignoring extra digits

- $|D_5D_4|D_3D_2|D_1S| < -|ZD_5|ZD_4|ZD_3|ZD_2|SD_1|$ where each 'Z' is a zone F
- PACK B(1),B(1) exchanges a byte's digits

<b>PACK</b>	PACK P(3),Z(4)						
•	I	<b>P(3</b> )	) <		Z(4	<b>ł</b> )	
Before:	??	??	??	F5	F4	F3	D2
After:	05	43	2D	F5	F4	F3	D2

PACK P(2),Z(4)
P(2) <---- Z(4)</p>

Before: ?? ?? F5 F4 F3 C2
After: 43 2C F5 F4 F3 C2

#### **label UNPK** $D_1(L_1, B_1), D_2(L_2, B_2)$

- The rightmost byte of the second operand is placed in the rightmost byte of the first operand, with zone (sign) and numeric digits reversed
- The remaining numeric digits from operand 2 are placed in the numeric digits of operand 1, and the zone digits of all but the rightmost byte of operand 1 are set to F, filling with X'F0' or ignoring extra digits

 $|ZD_{5}|ZD_{4}|ZD_{3}|ZD_{2}|SD_{1}| < -|D_{5}D_{4}|D_{3}D_{2}|D_{1}S|$ where each 'Z' is a zone F

**UNPK B(1),B(1)** exchanges a byte's digits

UNPK	Z(5)	,P(3)
------	------	-------

- **Z(5)** <---- P(3)
- Before: ?? ?? ?? ?? ?? 12 34 5C
- After: F1 F2 F3 F4 C5 12 34 5C

UNPK Z(4),P(2)

- Z(4) <---- P(2)</li>
   Before: ?? ?? ?? 12 3F
- After: F0 F1 F2 F3 12 3F

- These two RX instructions provide conversions between packed decimal and binary formats
- With PACK and UNPK, we can now convert between zoned and binary formats

#### **|** label CVB $R_1, D_2(X_2, B_2)$

Causes the contents of R<sub>1</sub> to be replaced by the binary representation of the <u>packed</u> decimal number in the doubleword (on a doubleword boundary) addressed by operand 2

A data exception (0007) occurs if operand 2 is not a valid packed decimal number

A fixed-point divide exception (0009) occurs if the result is too large to fit in a 32-bit word

For example:

	CVB	3,Z
	• • •	
Z	DS	0D
	DC	PL8'-2'

will convert 000000000000002D at location Z (data type D has doubleword alignment) to FFFFFFE in register 3

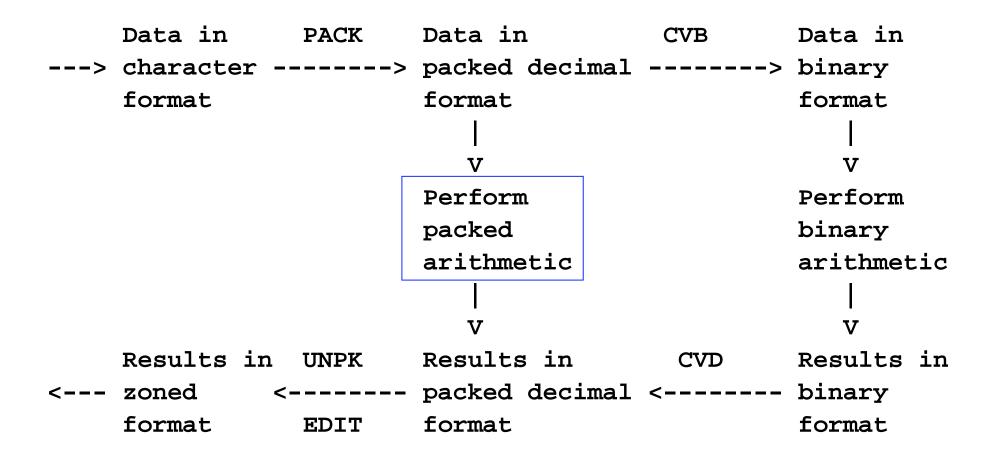
#### $= label CVD R_1, D_2(X_2, B_2)$

Causes the contents of the doubleword (on a doubleword boundary) addressed by operand 2 to be replaced by the packed decimal representation of the binary number in R<sub>1</sub>

Note that the "data movement" is left to right (like ST)

The exceptions which apply to CVB (0007 and 0009) do not apply to CVD

#### **Numeric Data Conversion Summary**



Getting results in nice character format, instead of just zoned, requires use of EDIT instruction

#### **Decimal Arithmetic**

- The box encloses the only subject on the previous slide which remains to be addressed: decimal arithmetic
- There isn't enough time to cover the decimal arithmetic instructions in detail, but they all have the following characteristics

#### **Decimal Arithmetic**

Two memory operands, each with its own length

Condition code is set similar to binary equivalents

In almost all cases (except operand 1 in ZAP), the operands must be valid packed decimal numbers, else an interrupt 0007 occurs (very popular!)

#### **Decimal Arithmetic**

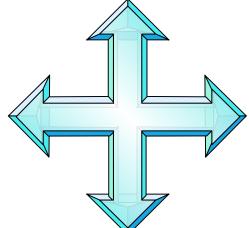
Here are the available instructions

- AP ADD DECIMAL
- CP COMPARE DECIMAL
- DP DIVIDE DECIMAL
- MP MULTIPLY DECIMAL
- SRP SHIFT AND ROUND DECIMAL
- SP SUBTRACT DECIMAL
- ZAP ZERO AND ADD DECIMAL

With the possible exception of SRP, these are easy to understand - see PoO

# **Instructions for Logical Operations**

# To Which We Must Say Yes or No



- Consider the four possible combinations of 2 bits
  - a = 0 0 1 1b = 0 1 0 1

These lead to the following binary relations

- a AND b = 0 0 0 1
- **a** OR b = 0 1 1 1
- **a** XOR b = 0 1 1 0

And these relations lead to the following twelve new instructions:

	RR Format	RX Format	SI Format	SS Format
AND Operation	NR	N	NI	NC
OR Operation	OR	Ο	OI	OC
XOR Operation	XR	x	XI	XC

anything with	itself	zero	one
AND	It remains unchanged	It is changed to zero	It remains unchanged
OR	It remains unchanged	It remains unchanged	It is changed to one
XOR	It is changed to zero	It remains unchanged	It is inverted

- All twelve instructions set the condition code:
  - 0 Result is zero
  - 1 Result is not zero

As an example, to change a zoned decimal number to EBCDIC, we have to force the rightmost zone to be F instead of a sign; so, if ZNUM is a three-byte zoned number, the following instruction will work: (why?)

OI ZNUM+2,X'F0'

To zero a register, we normally use SR, but a faster way to zero R5 (for example) is
 XR 5,5

To set bit 0 of BYTE to 1 while leaving the other bits unchanged

OI BYTE,B'1000000'

To set bit 0 of BYTE to 0 while leaving the other bits unchanged

NI BYTE,B'01111111'

# To invert bit 0 of BYTE to 1 while leaving the other bits unchanged XI BYTE, B'1000000'

To round the address in R7 down to the previous fullword boundary

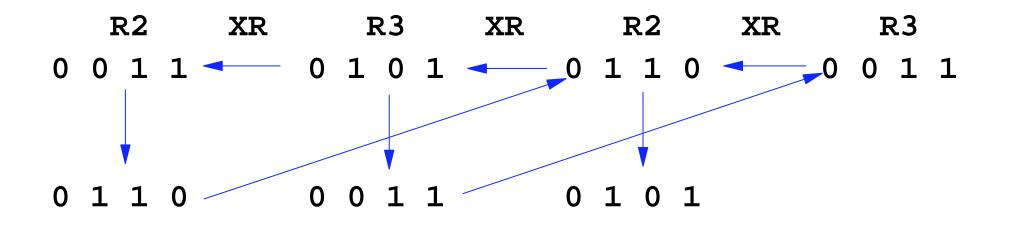
N 7,=X'FFFFFFC'

To round it up to the next fullword boundary

- LA 7,3(,7)
- N 7,=X'FFFFFFC'

- To exchange the contents of two registers without using any temporary space, use XR three times, alternating registers
- Memory contents can be exchanged similarly by using XC instead of XR
  - XR 2,3 Exchange
  - XR 3,2 contents of
  - XR 2,3 registers 2 and 3

Here's how that works, demonstrated using "mini" registers with all four posible bit combinations



# **The TEST UNDER MASK Instruction**

#### **] label** TM $D_1(B_1), I_2$

TM sets the condition code to reflect the value of the tested bits (those corresponding to 1-bits in the I<sub>2</sub> operand)

- O Selected bits all zeros, or the I<sub>2</sub> mask was zero
- 1 Selected bits mixed zeros and ones
- 2 --- (not set)
- 3 Selected bits all ones

## **The TEST UNDER MASK Instruction**

- Note that after TM, the extended branch mnemonics are interpreted as
  - BZ Branch if tested bits are Zeros, or mask is zero
  - BM Branch if tested bits are Mixed zeros and ones
  - BO <u>Branch if tested bits are Ones</u>

## **The TEST UNDER MASK Instruction**

- To determine if the first bit of BYTE is one
   TM BYTE,B'1000000'
- To check if BYTE is zero or blank (X'00' or X'40')
  - TM BYTE,B'10111111'
     BZ BLKZRO



In Which We Learn That Only a Small Fraction of the Assembler Language Has Been Covered

#### **Summary**

Five hours is just a start, but a good one

The one-semester course at NIU has

- More than 35 hours of lecture
- A dozen programs (almost one each week)
- Three exams

 The second course is Data Structures, and all program assignments are in assembler
 This is good reinforcement

#### What Wasn't Covered

Shift instructions, logical and arithmetic

Frequently used, but difficult instructions

- Edit (ED) and Edit and Mark (EDMK)
- Execute (EX)
- Translate (TR) and Translate and Test (TRT)

#### Floating point instructions

- Hexadecimal (the original)
- Binary (IEEE standard, recently added)

#### What Wasn't Covered

- Many general instructions added over the past twenty-five years, such as
  - Relative BRANCH instructions (no base register needed)
  - Instructions which reference a halfword (immediate) operand within the instruction
  - Instructions to save and set the addressing mode (24-bit or 31-bit)
  - And, most recently, the z/Architecture instructions to deal with 64-bit registers and addresses

#### What Wasn't Covered

- Privileged instructions
- The macro language, including conditional assembly (also available outside macros)
- The USING instruction, extended to allow implicit addresses everywhere
- External subroutines and register save area linkage conventions

#### Nevertheless...

You now have a basic understanding of S/390 and z/Architecture

You have seen what comprises a program written in assembler language

And you are ready, if you wish, to begin writing programs and go on to the next step



# **Congratulations!**

